

# 1 Pullbacks

**Definition 1.1.** Suppose we had morphisms  $f : A \rightarrow C$  and  $g : B \rightarrow C$ . Then the pullback of  $f$  and  $g$  is a pair of arrows  $p_1 : P \rightarrow A$  and  $p_2 : P \rightarrow B$  such that  $fp_1 = gp_2$ , and given any  $z_1 : Z \rightarrow A$  and  $z_2 : Z \rightarrow B$  such that  $z_1 = p_1u, z_2 = p_2u$ , there is a unique  $u : Z \rightarrow P$  such that  $z_1 = p_1u$  and  $z_2 = p_2u$

Pullbacks are unique, so we can denote the pullback of  $C$  as  $A \times_C B$

**Proposition 1.2.** Suppose  $C$  has products and equalizers. Suppose we had morphisms  $f : A \rightarrow C$  and  $g : B \rightarrow C$ . Then if  $A \times B$  is the product of  $A$  and  $B$ , and  $E$  is an equalizer of  $f\pi_1$  and  $g\pi_2$ , where  $p_1 = \pi_1e, p_2 = \pi_2e$ ,  $p_1, p_2$  is a pullback of  $f, g$ .

**Lemma 1.3** (Two pullback lemma). This diagram takes a long time to draw. Basically, if the 2 squares are pullbacks, the outer rectangle is a pullback. If the right square and outer rectangle are pullbacks, then the left square is too

**Corollary 1.4.** The pullback of a commutative triangle is a commutative triangle

# 2 Limits

**Definition 2.1.** Let  $J, C$  be categories. Then a diagram of type  $J$  in  $C$  is a functor  $D : J \rightarrow C$ .

**Definition 2.2** (Objects in the cone category). A cone to a diagram  $D$  is an object,  $C \in C$  and a collection of morphisms,  $\{c_j\}$ ,  $c_j : C \rightarrow D_j$  such that for every  $\alpha : i \rightarrow j$  in  $J$ ,  $D_\alpha c_i = c_j$ .

**Definition 2.3** (Morphisms in the cone category). A morphism of cones  $\vartheta : (C, \{c_j\}) \rightarrow (C', \{c'_j\})$  is a morphism  $\bar{\vartheta} : C \rightarrow C'$  in  $C$  such that for every  $j \in J$  we have  $c_j = c'_j \bar{\vartheta}$ .

**Definition 2.4.** If  $D : J \rightarrow C$  is a diagram, then a limit for  $D$  is a terminal object in the category of cones to  $D$ . If  $J$  is finite then the limit is called a finite limit.

The limit object would be denoted  $\lim_{\leftarrow j} D_j$ . It of course comes with a family of morphisms  $\{p_i\}$  such that  $p_i : \lim_{\leftarrow j} D_j \rightarrow D_i$ . This object has the property that for any cone  $(C, \{c_j\})$  to  $D$ , there is a unique  $u : C \rightarrow \lim_{\leftarrow j} D_j$  such that for every  $j \in J$  we have  $p_j \circ u = c_j$ .

We can now view products as a limit. Let  $\mathbf{J}$  be the discrete category with 2 objects, 2 morphisms (which both have to be identities). Then  $D : \mathbf{J} \rightarrow \mathbf{C}$  is a pair of objects  $D_1, D_2 \in \mathbf{C}$ . A cone of  $D$  is a object  $C \in \mathbf{C}$  together with morphisms  $c_i : C \rightarrow D_i$ . A limit of  $D$  would be a terminal cone, but this exactly coincides with the product.

Now we can construct equalizers with limits. Let  $\mathbf{J}$  be the category with 2 objects, 1, 2 and morphisms  $\alpha, \beta : 1 \rightarrow 2$  (of course the objects would need identity morphisms too but ignore those). Then a diagram of type  $\mathbf{J}$  would be 2 objects:  $D_1, D_2$  and morphisms  $D_\alpha, D_\beta : D_1 \rightarrow D_2$ . A cone would be  $c_i : C \rightarrow D_i$  such that  $D_\alpha c_1 = c_2 = D_\beta c_1$ , so a limit for  $D$  would be an equalizer of  $D_\alpha, D_\beta$ .

# 3 Continuity

**Definition 3.1.** Let  $F : \mathbf{C} \rightarrow \mathbf{D}$  be a functor. Then  $F$  preserves limits of type  $J$  if given a diagram  $D : \mathbf{J} \rightarrow \mathbf{C}$  and a limit  $p_j : L \rightarrow D_j$  then the cone  $F(p_j) : F(L) \rightarrow F(D_j)$  is a limit for the diagram  $F(D) : \mathbf{J} \rightarrow \mathbf{D}$ . If  $F$  preserves all limits, it is continuous.

**Proposition 3.2.** Let  $C$  be a locally small category. Then the representable functors  $\text{Hom}(C, -)$  is continuous.

Weird that this definition is being introduced now but not in duality but:

**Definition 3.3.** Let  $\mathbf{C}, \mathbf{D}$  be categories. Then  $F : \mathbf{C}^{\text{op}} \rightarrow \mathbf{D}$  is a contravariant functor on  $\mathbf{C}$ , where if  $f : A \rightarrow B$  is a morphism in  $\mathbf{C}$  then it is mapped to  $F(f) : F(B) \rightarrow F(A)$ , and  $F(g \circ f) = F(f) \circ F(g)$

Honestly I did not really understand the colimit part

**Definition 3.4.** A pushout is a pullback where you flip all the arrows and that's what it is and I don't want to draw the commutative diagram because it takes too long

**Example 3.5.** Let  $S^1, D^2$  be as defined in topology. Then  $S^2$  is the pushout of 2 of the same inclusion map  $i : S^1 \rightarrow D^2$ . //